## Geometric combinatorics of paths and deformations of convex polytopes



19 October 2023

- $\rightarrow$  Mon site : "Germain Poullot" dans Google
- $\rightarrow$  Onglet "Petit jeu"
- $\rightarrow$  Suivez les indications, mot de passe : 0000
- $\Rightarrow$  Amusez-vous !

Only in French, sorry ...

Merci Guillaume !!!

## Geometric combinatorics of paths and deformations of convex polytopes



19 October 2023





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#### 1 What is "Combinatorics of Polytopes"?

#### 2 Generalized permutahedra

- Deformations
- Submodular Cone
- Ongoing work

#### 3 Max-slope Pivot Polytopes

- Max-slope pivot rule
- Poset of slopes
- Pivot rule polytope of products of simplices

## What is "Combinatorics of Polytopes"?

*Polytope*: convex hull of finitely many points in  $\mathbb{R}^n$ 























## Representing polytopes















Dodecahedron the Universe



Icosahedron Water

## Representing polytopes



Tetrahedron Fire



Dodecahedron the Universe



Icosahedron Water



## Representing polytopes



Tetrahedron Fire

ø



Dodecahedron the Universe

ø



Hexahedron Earth



Octahedron



12









Water

2

ø





$$\textit{Face:} \ \mathsf{P}^{\textit{c}} := \left\{ \textit{\textit{x}} \in \mathbb{R}^{n} \ ; \ \left< \textit{\textit{x}}, \textit{\textit{c}} \right> = \mathsf{max}_{\textit{\textit{y}} \in \mathsf{P}} \left< \textit{\textit{y}}, \textit{\textit{c}} \right> \right\}$$



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*One way*: Take a polytope  $\rightarrow$  combinatorial info (e.g. face lattice)

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$$\Pi_n = \operatorname{conv} \left\{ \begin{pmatrix} \sigma(1) \\ \vdots \\ \sigma(n) \end{pmatrix} ; \ \sigma \text{ permutation of } \{1, \dots, n\} \right\}$$

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Germain Poullot

## Generalized permutahedra
Coarsening: Choose maximal cones and merge them

### Definition

Q is a *deformation* of P iff  $\mathcal{N}_Q$  coarsens  $\mathcal{N}_P$ .



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*Braid fan*: arrangement of hyperplanes  $H_{i,j} := \{ \mathbf{x} ; x_i = x_j \}$ 

# Braid fan

### Definition

Braid fan: arrangement of hyperplanes  $H_{i,j} := \{ \mathbf{x} ; x_i = x_j \}$ 

### Definition

Generalized permutahedron: deformation of  $\Pi_n$ 

i.e. P generalized permutahedron iff  $\mathcal{N}_{\mathsf{P}}$  coarsens braid fan





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i.e. P generalized permutahedron iff  $\mathcal{N}_{\mathsf{P}}$  coarsens braid fan



 $\mathcal{P}(\mathsf{P})$ : all the posets associated to faces of  $\mathsf{P}$ 

# Deformations of $\Pi_4$



Sequence of deformations of  $\Pi_4$ 

Minkowski sum:  $P + Q = \{ \boldsymbol{p} + \boldsymbol{q} ; \boldsymbol{p} \in P, \boldsymbol{q} \in Q \}$ 

#### Theorem

If Q, R deformations of P, then:

for all  $\lambda > 0$ ,  $\lambda Q$  deform. of P Q + R deform. of P

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### Definition

Deformation cone:  $\mathbb{DC}(P) := \{Q ; Q \text{ deformation of } P\}$  is a cone.



Parametrization:

height vector:  $\boldsymbol{h} = (h_r)_{r \text{ rays}}$ 

Minkowski sum:  $P + Q = \{ \boldsymbol{p} + \boldsymbol{q} ; \boldsymbol{p} \in P, \boldsymbol{q} \in Q \}$ 

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## Parametrization:

*Submodular cone*: deformation cone of the permutahedron  $\Pi_n$ 

	$\mathbb{DC}(\Pi_n)$	
Dim (no lineal)	$2^{n} - n - 1$	
# facets	$\binom{n}{2}2^{n-2}$	
# rays	unknown!	

# Submodular Cone for $\Pi_3$



Submodular cone: deformation cone of the permutahedron  $\Pi_n$ 

	$\mathbb{DC}(\Pi_n)$	
Dim (no lineal)	$2^{n} - n - 1$	
# facets	$\binom{n}{2}2^{n-2}$	
# rays	unknown!	

Submodular cone: deformation cone of the permutahedron  $\Pi_n$ 

## Theorem (Faces of $\mathbb{DC}(\mathsf{P})$ )

If Q deformation of P, then  $\mathbb{DC}(Q)$  is a face of  $\mathbb{DC}(P)$ .

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If Q deformation of P, then  $\mathbb{DC}(Q)$  is a face of  $\mathbb{DC}(P)$ .

	$\mathbb{DC}(\Pi_n)$	$\mathbb{DC}(Asso_n)$	
Dim (no lineal)	$2^{n} - n - 1$	$\binom{n}{2}$	
# facets	$\binom{n}{2}2^{n-2}$	$\binom{\overline{n}}{2}$	
# rays	unknown!	$\binom{\overline{n}}{2}$	
		is simplicial!	

Submodular cone: deformation cone of the permutahedron  $\Pi_n$ 

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If Q deformation of P, then  $\mathbb{DC}(Q)$  is a face of  $\mathbb{DC}(P)$ .

	$\mathbb{DC}(\Pi_n)$	$\mathbb{DC}(Asso_n)$	$\mathbb{DC}(Z_G)$	$\mathbb{DC}(N_B)$
Dim (no lineal)	$2^{n} - n - 1$	$\binom{n}{2}$	Ν	N
# facets	$\binom{n}{2}2^{n-2}$	$\binom{\overline{n}}{2}$	Е	E
# rays	unknown!	$\binom{\overline{n}}{2}$	Х	Х
		is simplicial!	Т	Т

$$G = (V, E)$$
 a graph,  $n = |V|$ 

### Definition

Graphical zonotope 
$$Z_G := \sum_{(i,j) \in E} [e_i, e_j]$$

 $Z_G$  deformation of  $\Pi_n \implies \mathbb{DC}(Z_G)$  is a face of  $\mathbb{DC}(\Pi_n)$ 





## Theorem (Padrol, Pilaud, P., '23)

Explicit facet-description of  $\mathbb{DC}(Z_G)$ 

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### Corollary

 $dim \mathbb{DC}(Z_G) = \# cliques of G$ # facets of  $\mathbb{DC}(Z_G) = \sum_{(i,j)\in E} 2^{|\{k : (i,k), (j,k)\in E\}|}$ 

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### Corollary

 $\mathbb{DC}(Z_G)$  simplicial iff G without triangle

**NB**: Recover facet-description of  $\mathbb{DC}(\Pi_n)$ 



## Definition

Building set  $B \subseteq 2^{[n]}$  with:  $X_{1,2} \in B, X_1 \cap X_2 \neq \emptyset \Rightarrow X_1 \cup X_2 \in B$ 

### Definition

*Nestohedron*  $N_B := \sum_{X \in B} \Delta_X$  where  $\Delta_X = \text{conv}\{e_i ; i \in X\}$ 

 $N_B$  deformation of  $\Pi_n \implies \mathbb{DC}(N_B)$  is a face of  $\mathbb{DC}(\Pi_n)$ 



Elementary blocks  $X \in \varepsilon(B)$  iff X is not a union Maximal block  $\mu(X) := \max\{Y \in B ; Y \subsetneq X\}$ 

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### Corollary

dim  $\mathbb{DC}(N_B) = |B| - \#$  singletons # facets of  $\mathbb{DC}(N_B) = |\varepsilon(B)| + \sum_{X \in B \setminus \varepsilon(B)} {|\mu(X)| \choose 2}$ 

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### Corollary

 $\mathbb{DC}(N_B)$  simplicial iff B has no non-elementary block with 3 maximal subblocks

**NB**: Recover facet-description of  $\mathbb{DC}(\Pi_n)$ 



# Ongoing work - Hypergraphic polytopes

## Definition

Hypergraphic pol  $\mathsf{P}_H := \sum_{X \in H} \Delta_X$  with  $H \subseteq 2^{[n]}$ 


### Ongoing work - Quotientopes

#### Definition

Quotientopes: Minkowski sum of shard polytopes



# Max-slope Pivot Polytopes

#### Linear optimization



### Simplex method











Linear optimization in dimension 2 (simplex method): EASY !



Convention: choose upper











Optimization in higher dimension: make it 2-dimensional !



Max-slope pivot rule: take (improving) neighbor with best slope

Optimization in higher dimension: make it 2-dimensional !



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 $\omega^{-}$ 

3








































## Monotone path polytope



*Coherent monotone path*: path obtained via max-slope pivot rule

Monotone path fan: Fan with  $\omega \sim \omega'$  iff same path

# Monotone path polytope



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#### Theorem (Billera, Sturmfels, '92)

The monotone path fan is polytopal. Monotone path polytope  $\Sigma_c(P)$ : dual to monotone path fan

# Monotone path polytope



*Coherent monotone path*: path obtained via max-slope pivot rule

 $\begin{array}{l} \textit{Monotone path fan: Fan with} \\ \boldsymbol{\omega} \sim \boldsymbol{\omega}' \textit{ iff same path} \end{array}$ 

#### Theorem (Billera, Sturmfels, '92)

The monotone path fan is polytopal. Monotone path polytope  $\Sigma_c(P)$ : dual to monotone path fan

$$\Sigma_{c}(\Delta_{d}) = \mathsf{Cube}_{d-1} \qquad \Sigma_{c}(\mathsf{Cube}_{d}) = \Pi_{d}$$


















































































*Coherent arborescence*: arborescence obtained via max-slope pivot rule

#### Pivot rule fan:

 $oldsymbol{\omega}\simoldsymbol{\omega}'$  iff same arborescence.



Coherent arborescence: arborescence obtained via max-slope pivot rule Pivot rule fan:  $\omega \sim \omega'$  iff same arborescence. Pivot rule fan refines monotone path fan



#### Theorem (Black, De Loeara, Lütjeharms, Sanyal '22)

The pivot rule fan is polytopal.

(Max-slope) pivot polytope  $\Pi_c$ : dual to the pivot rule fan



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The pivot rule fan is polytopal.

(Max-slope) pivot polytope  $\Pi_c$ : dual to the pivot rule fan

$$\Sigma_{\boldsymbol{c}}(\Delta_d) = \mathsf{Cube}_{d-1}$$

 $\Pi_{c}(\Delta_{d}) = \operatorname{Asso}_{d}$ 



 $\Pi_{\boldsymbol{c}}(\Delta_d) \simeq \mathsf{Asso}_d$ 

for all **c** 





 $\Pi_{\boldsymbol{c}}(\mathsf{Cube}_d) \simeq \Pi_d \qquad \qquad \text{for all } \boldsymbol{c}$ 

Conjecture (Pilaud, Sanyal)

 ${\sf \Pi}_{m{c}}({\sf \Delta}_{d_1} imes {\sf \Delta}_{d_2})\simeq {\sf Asso}_{d_1}\star {\sf Asso}_{d_2}$ 

 $\star$  shuffle product



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Challenge 1: Prove the conjecture!



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Challenge 1: Prove the conjecture! Challenge 2: Give **geometric** proofs!





Max-slope pivot rule: take (improving) neighbor with best slope

For  $\omega$ , what is important?



Max-slope pivot rule: take (improving) neighbor with best slope

#### For $\omega$ , what is important?

Slopes: 
$$au_{\omega}(u,v) = rac{\langle \omega,u-v
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 $\theta: \mathbb{R}^d \to \mathbb{R}^m$ , injective linear map



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What is **really** important?? The comparisons of slopes! Compare coordinates of  $\theta(\omega)$ 

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 $\theta : \mathbb{R}^d \to \mathbb{R}^m$ , injective linear map

What is **really** important?? The comparisons of slopes! Compare coordinates of  $\theta(\omega)$ Where is  $\theta(\omega)$  in the braid fan  $\mathcal{B}_m$ ?



Too many edges



Too many edges, **but** parallelism saves us!



Too many edges, **but** parallelism saves us!

Geometric proof of  $\Pi_c(\text{Cube}_d) = \Pi_d$ 





Too many edges



Too many edges, but affine independence saves us!



Too many edges, **but** affine independence saves us! Geometric proof of  $\Pi_c(\Delta_d) = \text{Asso}_d$ 



### Shuffles

```
Shuffle: (E, \leq) and (F, \preceq) posets, then \trianglelefteq is a shuffle when:
groud set : E \sqcup F
relations : all relations of \leq ; all relations of \preceq ;
for each e \in E, f \in F, choose if e \trianglelefteq f or e \trianglerighteq f
(+ \text{ transitive closure})
```

### Shuffles

### Shuffle: $(E, \leq)$ and $(F, \preceq)$ posets, then $\trianglelefteq$ is a shuffle when: groud set : $E \sqcup F$ relations : all relations of $\leq$ ; all relations of $\preceq$ ; for each $e \in E$ , $f \in F$ , choose if $e \trianglelefteq f$ or $e \trianglerighteq f$ (+ transitive closure)

#### Theorem (Chapoton, Pilaud '22)

P, Q: generalized permutahedra. Exists polytope P \* Q s.t.

 $\mathcal{P}(\mathsf{P} \star \mathsf{Q}) = \{ all \ shuffles \ between \ \leq \in \mathcal{P}(\mathsf{P}) \ and \ \leq \in \mathcal{P}(\mathsf{Q}) \}$ 

Combine parallelism & affine independence:



### Combine parallelism & affine independence:

TheoremFor 
$$\Delta_{d_1} \times \cdots \times \Delta_{d_r}$$
, all (generic) direction: $\Pi_{\boldsymbol{c}}(\Delta_{d_1} \times \cdots \times \Delta_{d_r}) \simeq \operatorname{Asso}_{d_1} \star \cdots \star \operatorname{Asso}_{d_r}$ 

### Example

(a) 
$$\Pi_{\boldsymbol{c}}(\Box_d) \simeq \operatorname{Perm}_d$$
  
(b)  $\Pi_{\boldsymbol{c}}(\Box_m \times \Delta_n) \simeq (m, n)$ -multiplihedron  
(c)  $\Pi_{\boldsymbol{c}}(\Delta_m \times \Delta_n) \simeq (m, n)$ -constrainahedror
- 1) For which P,  $\Pi_c(P)$  is a generalized permutahedron?  $\longrightarrow$  a priori, only products of simplices, but no proof
- 2) Is  $\Pi_{c}(\mathsf{P})$  projection of a generalized permutahedron?  $\longrightarrow$  pivot fan sent inside  $\operatorname{Im}(\theta) \cap \mathcal{B}_{m}$
- 3) When  $\Pi_{c}(P)$  and  $\Pi_{c}(Q)$  **not** generalized permutahedra, what happen to  $\Pi_{c}(P \times Q)$ ?
- $\longrightarrow$  not equivalent to  $\Pi_{c}(\mathsf{P}) \star \Pi_{c}(\mathsf{Q})$ , but "embeds" in it

# What I have presented

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A A Vandermonde-like determinant

# Thank you!



Notations:  $Sx = S \cup \{x\}$ ,  $(f_X)_{X \subseteq [n]}$  canonical basis of  $\mathbb{R}^{2^{[n]}}$ 

#### Definition

Submodular vector 
$$\mathbf{n}(S, u, v) = \mathbf{f}_{Suv} - \mathbf{f}_{Su} - \mathbf{f}_{Sv} + \mathbf{f}_{S}$$
  
for  $u, v \in S \subseteq [n]$ 



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#### Lemma (Submodular normal)

n(S, u, v) are the facet's normals of  $\mathbb{DC}(\Pi_n)$ 

#### Lemma (Cubic relation)

 $u, v, x \notin S \subseteq [n]$ n(Suvx, u, v) + n(Sux, u, x) = n(Suv, u, v) + n(Suvx, u, x)



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# Monotone path polytope and pivot rule polytope

Let  $\mathsf{P} \subset \mathbb{R}^d$  be a polytope.

Max-slope pivot rule: 
$$A^{\omega}(v) = \operatorname{argmax} \left\{ \frac{\langle \omega, u-v \rangle}{\langle c, u-v \rangle}; u \text{ impr. neig. of } v \right\}.$$

*Coherent monotone path*: A monotone path that can be obtained via the max-slope pivot rule.

Monotone path polytope  $\Sigma_c(P)$  [?]: Fiber polytope of  $P \xrightarrow{\pi} Q$  with Q a segment. (Can be seen as a Minkowski sum of sections of P.) The vertices of  $\Sigma_c(P)$  are all coherent monotone paths.

*Coherent arborescence*: An arborescence that can be obtained via the max-slope pivot rule.

*Pivot rule polytope*  $\Pi_{c}(P)$ : Polytope which vertices are all coherent arborescences.

$$\Pi_{\boldsymbol{c}}(\mathsf{P}) = \operatorname{conv}\left\{\sum_{\nu \neq v_{opt}} \frac{1}{\langle \boldsymbol{c}, \mathcal{A}(\nu) - \nu \rangle} (\mathcal{A}(\nu) - \nu); \mathcal{A} \text{ coherent arbo. of } \mathsf{P}\right\}$$

# Monotone path polytope and pivot rule polytope

*Coherent arborescence*: An arborescence that can be obtained via the max-slope pivot rule.

*Pivot rule polytope*  $\Pi_{c}(P)$ : Polytope which vertices are all coherent arborescences. Can also be seen as a Minkowski sum of sections:

 $\sum_{v \in V(P)}$ (section between v and its improving neighbors)



Idea 1:

Fix a polytope P, and direction c, n vertices, m edges.

 $\theta : \mathbb{R}^d \to \mathbb{R}^m$  sends the pivot fan inside  $\operatorname{Im}(\theta) \cap \mathcal{B}_m$ Problem: This is not a braid fan as  $d \ll m...$ 

If m' classes of parallelism:  $\overline{\theta} : \mathbb{R}^d \to \mathbb{R}^{m'}$  sends the pivot fan inside  $\operatorname{Im}(\theta) \cap \mathcal{B}_{m'}$ *Problem*: This is not a braid fan as  $d \ll m' \ll m...$ 

We need to go lower dimensional!

*Idea 2*: Fix a polytope P, direction *c*, <u>*n* vertices</u>, *m* edges.

Fix A arborescence:  $\vartheta_A(\omega) = (\tau_{\omega}(u, A(u)); u \text{ vertex})$ 



 $\vartheta_A$ : linear, injective,  $\mathbb{R}^d \to \mathbb{R}^{n-1}$ **but** if  $\omega$  does not capture A, then  $\vartheta_A(\omega)$  have no meaning... Adapted slope map:  $\vartheta(\omega) = \vartheta_{A^{\omega}}(\omega)$ i.e. take  $\omega$  and look at the slope of the edges it selects.

# Case of the *d*-simplex

 $d = n - 1 \iff \mathsf{P}$  is a simplex For  $\Delta_d: \vartheta: \mathbb{R}^d \to \mathbb{R}^d$  piece-wise linear, ker  $\vartheta = \{\mathbf{0}\} \Rightarrow$  bijection  $\vartheta$  sends the pivot fan of  $\Delta_d$  inside  $\mathcal{B}_d$ .

